

Entropic Hardness of Module-LWE from Module-NTRU

Corentin JEUDY

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Potential Candidates: NIST PQC Standardization

NIST **PQC standardization process** launched in 2016. Finalists announced in July 2020¹:

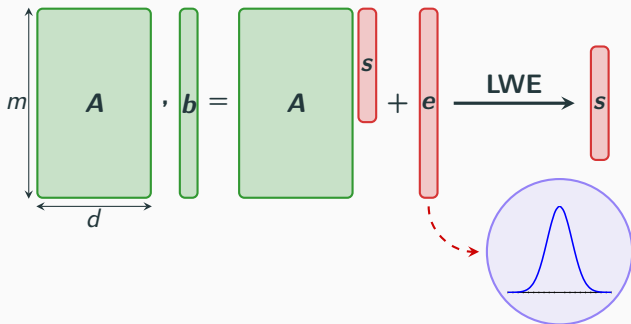
Encryption	Signature	
Crystals-Kyber Saber	Crystals-Dilithium	M-LWE & co
NTRU	Falcon	lattice-based
Classic McEliece	Rainbow ⚠	

Important to study the **hardness** of the underlying assumptions
e.g. **M-LWE**

¹Third round “winners” were supposed to be announced at the end of March. We must wait a bit longer.

Warm-Up: The Learning With Errors (LWE) Problem

The **Learning With Errors** problem was introduced in [Reg05]².



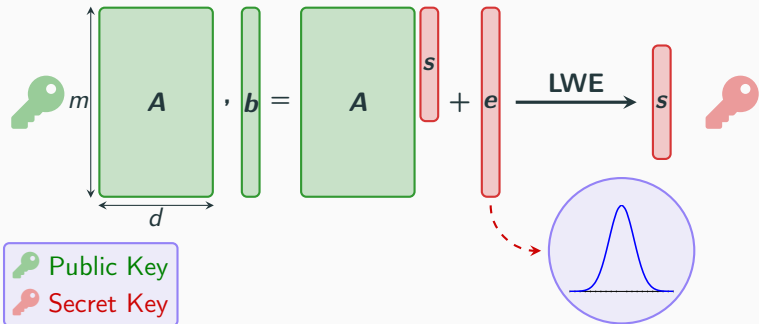
where $\mathbf{A} \leftarrow \mathcal{U}(\mathbb{Z}_q^{m \times d})$, $\mathbf{s} \leftarrow \mathcal{U}(\mathbb{Z}_q^d)$, and \mathbf{e} Gaussian.

LWE is proven to be at least as hard as hard problems on lattices.

²O. Regev, *On Lattices, Learning With Errors, Random Linear Codes, and Cryptography*, STOC'05

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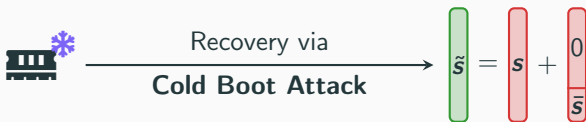
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Key Recovery for PKE is exactly the **LWE** problem

²O. Regev, *On Lattices, Learning With Errors, Random Linear Codes, and Cryptography*, STOC'05

Motivational Example: Key Recovery via Cold Boot Attack

1. Physical attack to recover a noisy secret \tilde{s} .



2. Target a new LWE instance with

$$\Delta b = b - A\tilde{s} = \begin{matrix} \text{[Green Box]} \\ A \end{matrix} \begin{matrix} 0 \\ \bar{s} \end{matrix} - e$$

?

Under what condition on \bar{s} is the problem still hard?
 \bar{s} must have enough **entropy** \rightarrow **Entropic hardness**

Adding an Algebraic Structure

Replace \mathbb{Z} with a ring $\mathcal{R} = \mathbb{Z}[x]/\langle f(x) \rangle$, e.g., $f(x) = x^n + 1$ with $n = 2^\ell$ and \mathbb{Z}_q by $\mathcal{R}_q = \mathbb{Z}_q[x]/\langle f(x) \rangle$

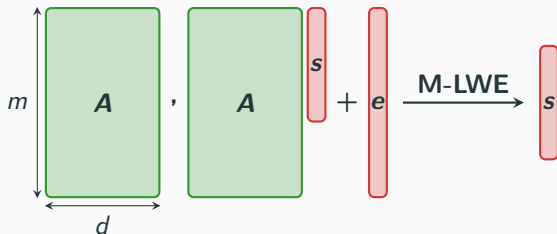
$$\sum_{i=0}^{n-1} a_i \cdot x^i \in \mathcal{R} \xleftrightarrow{\text{embedding}} \begin{bmatrix} a_0 \\ \vdots \\ a_{n-1} \end{bmatrix} \in \mathbb{Z}^n$$
$$\left(\sum_{i=0}^{n-1} a_i \cdot x^i \right) \cdot \left(\sum_{i=0}^{n-1} b_i \cdot x^i \right) \xleftrightarrow{\text{Rot}(\mathbf{a})} \begin{bmatrix} b_0 \\ \vdots \\ b_{n-1} \end{bmatrix}$$

The diagram illustrates the mapping of polynomial multiplication in the ring \mathcal{R} to a matrix-vector multiplication in \mathbb{Z}^n . The first part shows the embedding of a polynomial $\sum_{i=0}^{n-1} a_i \cdot x^i$ into a vector $\begin{bmatrix} a_0 \\ \vdots \\ a_{n-1} \end{bmatrix}$. The second part shows the multiplication of two polynomials $\left(\sum_{i=0}^{n-1} a_i \cdot x^i \right) \cdot \left(\sum_{i=0}^{n-1} b_i \cdot x^i \right)$ being mapped to a matrix-vector product where the matrix is a shaded box labeled $\text{Rot}(\mathbf{a})$ and the vector is $\begin{bmatrix} b_0 \\ \vdots \\ b_{n-1} \end{bmatrix}$.

Efficiency: FFT-like algorithms, use of structured matrices.

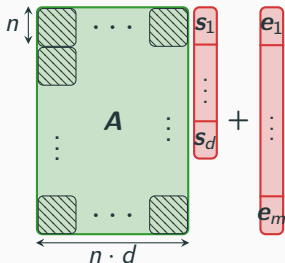
Storage: Structured matrices represented by a single vector.

Module-LWE as Structured-LWE

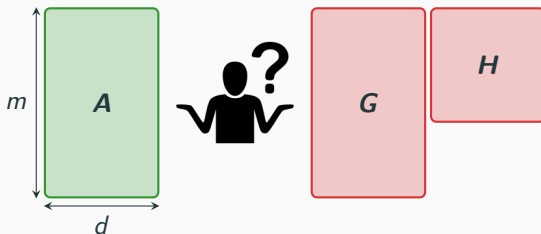


where $\mathbf{A} \leftarrow \mathcal{U}(\mathcal{R}_q^{m \times d})$, $\mathbf{s} \leftarrow \mathcal{U}(\mathcal{R}_q^d)$, and \mathbf{e} Gaussian.

It can be seen as a *Structured* LWE (**S-LWE**) with dimensions nm & nd .



What about Module-NTRU?



where $\mathbf{A} \leftarrow \mathcal{U}(\mathcal{R}_q^{m \times d})$, \mathbf{G}, \mathbf{F} Gaussian, and $\mathbf{H} = (\mathbf{F} \bmod q\mathcal{R})^{-1}$.

What do we know so far about Entropic Hardness?

LWE	M-LWE
[BD20a] ³ [BD20b] ⁵ (R-LWE)	[LWW20] ⁴ (ePrint) ★ Today

Our contribution:

★ M-LWE is hard with **arbitrary** s , if s has enough entropy.

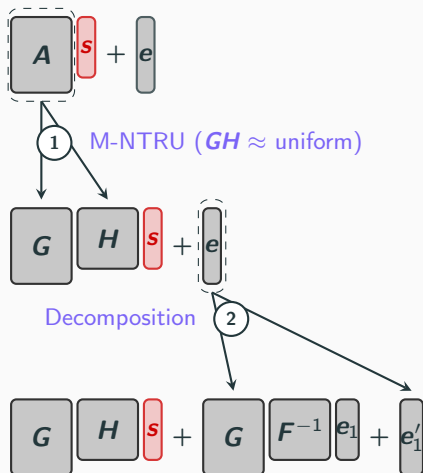
³Z. Brakerski, N. Döttling, *Hardness of LWE on General Entropic Distribution*, EUROCRYPT'20

⁴H. Lin, Y. Wang, M. Wang, *Hardness of Module-LWE and Ring-LWE on General Entropic Distribution*

⁵Z. Brakerski, N. Döttling, *Lossiness and Entropic Hardness for Ring-LWE*, TCC'20

Entropic Hardness of M-LWE

Replacing \mathbf{A} by \mathbf{GH} , with \mathbf{F} , \mathbf{G} Gaussian and \mathbf{H} the mod- q inverse of \mathbf{F} . The secret \mathbf{s} is only assumed to have **large enough entropy**. Based on the work by Brakerski and Döttling [BD20b] on R-LWE.



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$$\boxed{\mathbf{A}} \mathbf{s} + \mathbf{e}$$

① M-NTRU ($\mathbf{GH} \approx \text{uniform}$)

$$\boxed{\mathbf{G}} \boxed{\mathbf{H}} \mathbf{s} + \mathbf{e}$$

Decomposition

$$\boxed{\mathbf{G}} \boxed{\mathbf{H}} \mathbf{s} + \boxed{\mathbf{G}} \boxed{\mathbf{F}^{-1}} \mathbf{e}_1 + \mathbf{e}'_1$$

$$\boxed{\mathbf{G}} \left(\boxed{\mathbf{H}} \mathbf{s} + \boxed{\mathbf{F}^{-1}} \mathbf{e}_1 \right) + \mathbf{e}'_1$$

$$H_\infty(\mathbf{s} | \mathbf{G}(\mathbf{H}\mathbf{s} + \mathbf{F}^{-1}\mathbf{e}_1) + \mathbf{e}'_1)$$

$$\geq H_\infty(\mathbf{s} | \mathbf{H}\mathbf{s} + \mathbf{F}^{-1}\mathbf{e}_1)$$

$$\geq H_\infty(\mathbf{s} | \mathbf{H}\mathbf{s} + \mathbf{F}^{-1}\mathbf{e}_2) \quad (\mathbf{e}_2 \in \Lambda(\mathbf{F}))$$

$$= H_\infty(\mathbf{s} | \mathbf{s} + \mathbf{e}_2)$$

$$\geq H_\infty(\mathbf{s} | \mathbf{s} + \mathbf{e}') - nd \log_2 \|\mathbf{F}\|_2$$

$$\geq H_\infty(\mathbf{s}) - nd \log_2 \frac{q}{\sigma_{\mathbf{e}'}} - nd \log_2 \|\mathbf{F}\|_2$$

$$\textcircled{2} \quad \sigma_{\mathbf{e}} > \sigma_{\mathbf{e}'} \|\mathbf{GF}^{-1}\|_2$$

Our contribution

- ✓ Reduction from Module-NTRU to Module-LWE with **general⁶ secret distributions**.

Related Work

- 📄 Other reduction in [LWW20] from Module-LWE (uniform secret) to Module-LWE (general secret).
 - ✗ Not rank-preserving.
 - ✓ Assumption proven on module lattices.
 - = Parameter regimes with sometimes better or worse results.

Open Questions

- ? Reduction from module lattice problems to Module-NTRU?
- ? Prove the hardness of Module-LWE with low-entropy secret distributions without increasing the rank?

⁶with some restrictions though

Thank you for your
attention!



Questions?



Z. Brakerski and N. Döttling.

Hardness of LWE on general entropic distributions.

In *EUROCRYPT (2)*, volume 12106 of *Lecture Notes in Computer Science*, pages 551–575. Springer, 2020.



Z. Brakerski and N. Döttling.

Lossiness and entropic hardness for ring-lwe.

In *TCC (1)*, volume 12550 of *Lecture Notes in Computer Science*, pages 1–27. Springer, 2020.



H. Lin, Y. Wang, and M. Wang.

Hardness of module-lwe and ring-lwe on general entropic distributions.

IACR Cryptol. ePrint Arch., page 1238, 2020.



O. Regev.

On lattices, learning with errors, random linear codes, and cryptography.

In *STOC*, pages 84–93. ACM, 2005.

